

CS206 Data Structures

Recursion

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Class Objectives

- Understand the concept of recursion
- Can write a recursion function for addressing problems
 - Base case
 - Recursive calls
- Get to know that a program can run fast or slow depending on its algorithm

The Recursion Pattern

- Recursion: when a method calls itself
- Classic example: the factorial function:
 - $n! = 1 \cdot 2 \cdot 3 \cdot \dots \cdot (n-1) \cdot n$

- Recursive definition:

$$f(n) = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot f(n - 1) & \text{else} \end{cases}$$

- As a Java method:

```
// recursive factorial function
public static int recursiveFactorial(int n) {
    if (n == 0) return 1;    // basis case
    else return n * recursiveFactorial(n- 1);    // recursive case
}
```

Content of a Recursive Method

```
// recursive factorial function
public static int recursiveFactorial(int n) {
    if (n == 0) return 1;    // basis case
    else return n * recursiveFactorial(n- 1);    // recursive case
}
```

□ Base case(s)

- Return values w/o recursive calls; there should be at least one base case.

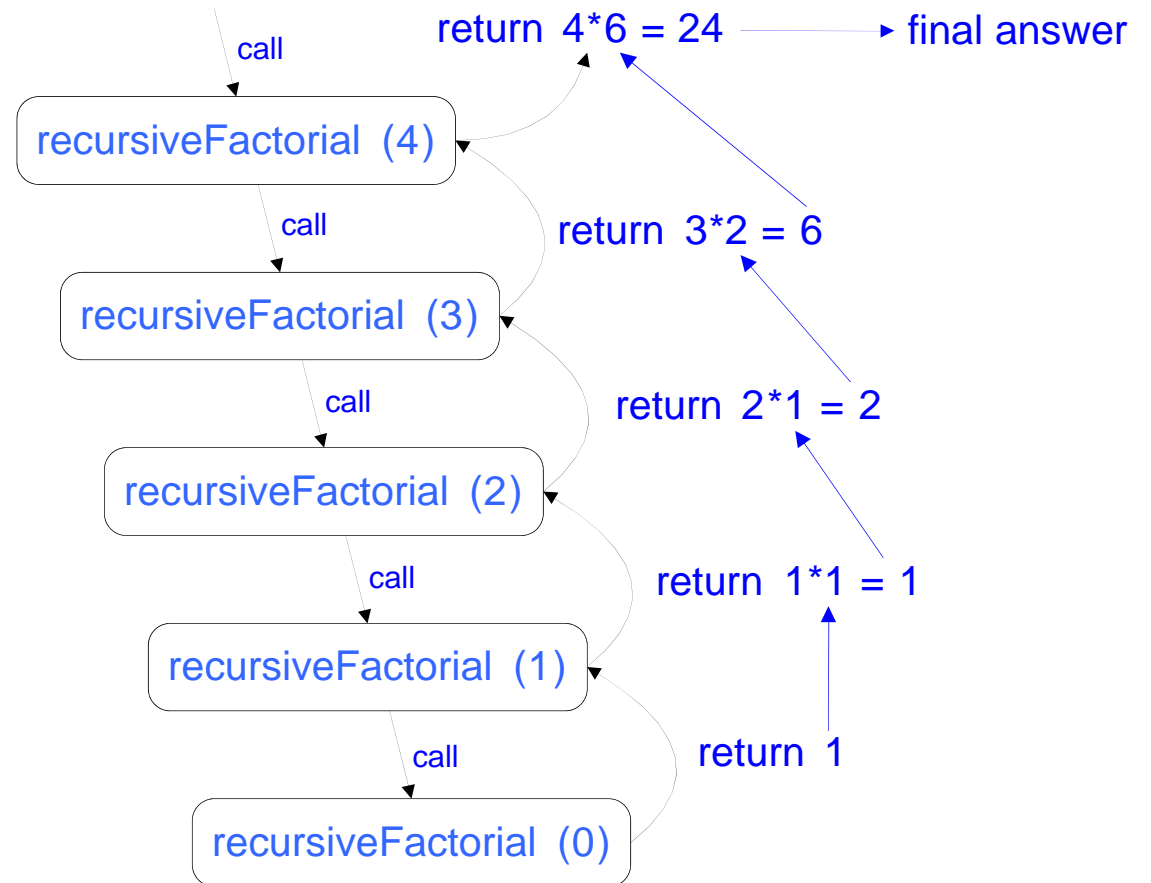
□ *Recursive calls*

- Calls to the current method that makes progress towards a base case.

Visualizing Recursion

- Recursion trace
- A box for each recursive call
- An arrow from each caller to callee
- An arrow from each callee to caller showing return value

Example recursion trace:



Linear Recursion

- Begin by testing for a set of base cases and recur once

Algorithm $\text{LinearSum}(A, n)$

Input: integer array A
and integer $n \geq 1$
s.t. A has at least n elements

Output: Sum of first n integers in A

if $n = 1$ then

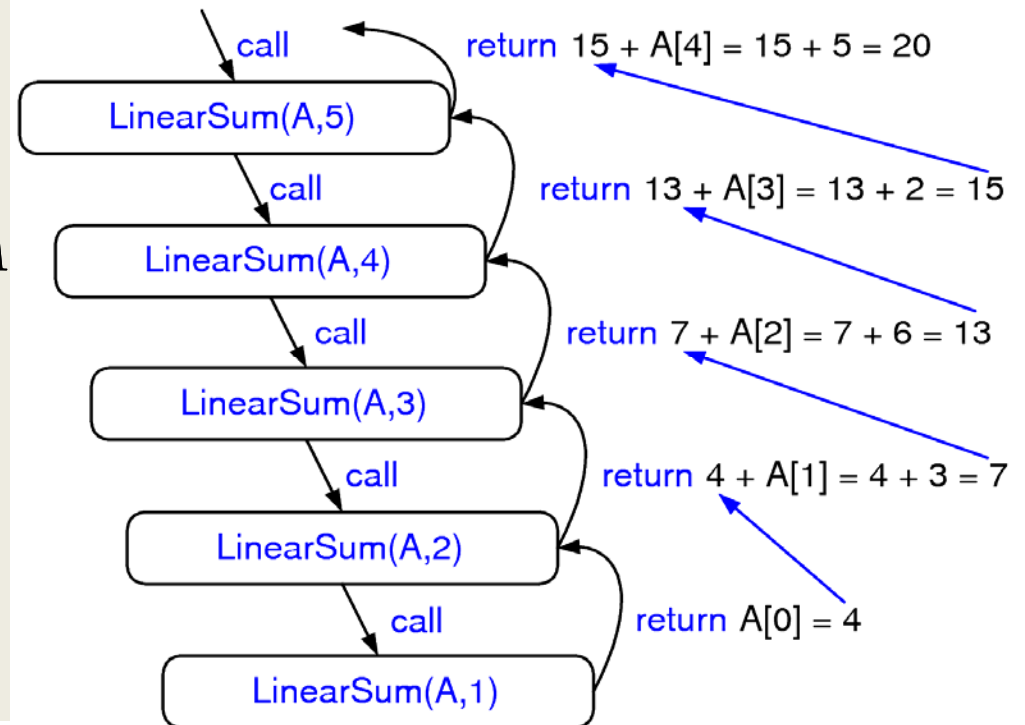
return $A[0]$

else

return $\text{LinearSum}(A, n - 1)$
+ $A[n - 1]$

end if

Example recursion trace:



- Iterative approach (using “for”) should be easier for this problem

- But this linear recursion is useful when the first (or last) plus a remaining set that has the same structure

Reversing an Array (HW: Self study)

Algorithm ReverseArray(A, i, j)

Input: array A and non-negative integer indices i and j

Output: Reversal of elements in A starting at index i and ending at j

if $i < j$ **then**

 Swap $A[i]$ and $A[j]$

 ReverseArray($A, i + 1, j - 1$)

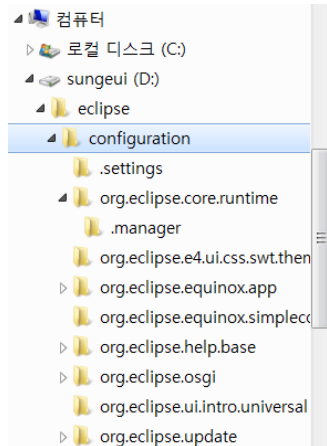
end if

return

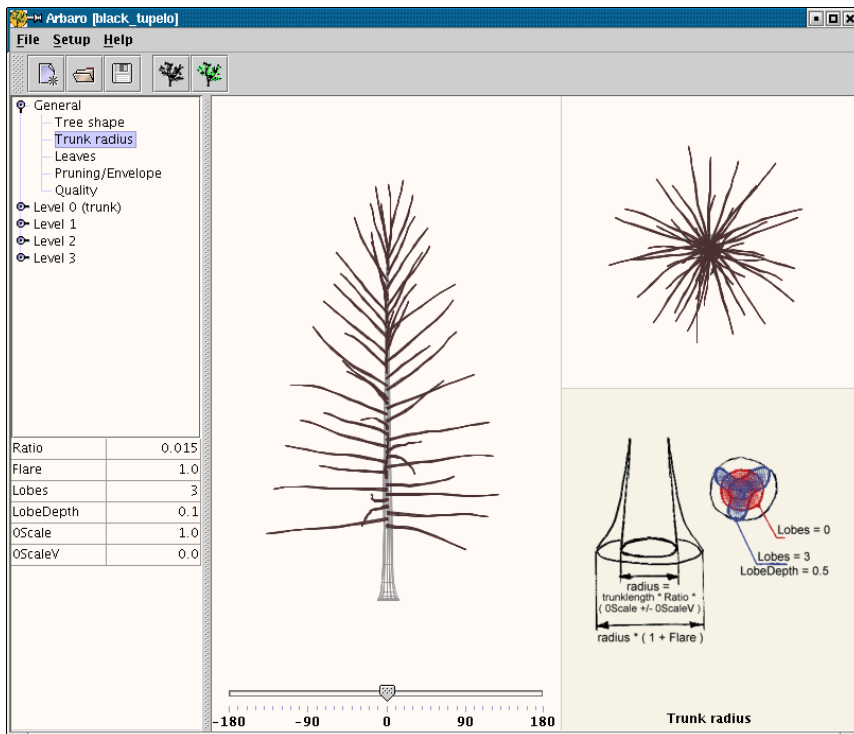
Defining Arguments for Recursion

- In creating recursive methods, it is important to define the methods in ways that facilitate recursion
- This sometimes requires to use additional parameters that are passed to the method
 - For example, we defined the array sum method as `LinearSum(A, n)`, not `LinearSum (A)`

Recursive structures are everywhere!



- Tree generation tool in computer graphics
 - Generate recursively



Google images

Computing Powers

□ The power function, $p(x, n) = x^n$, can be defined recursively:

$$p(x, n) = \begin{cases} 1 & \text{if } n = 0 \\ x \cdot p(x, n - 1) & \text{else} \end{cases}$$

□ This leads to a power function that invokes n times of multiplication (for we make n recursive calls).

- = runs linearly as a function of n

□ We can do better than this, however.

Recursive Squaring

□ We can derive a more efficient linearly recursive algorithm by using repeated squaring:

$$p(x, n) = \begin{cases} 1 & \text{if } n = 0 \\ p(x, \frac{n}{2})^2 & \text{if } n > 0 \text{ is even} \\ x \cdot p(x, \frac{n-1}{2})^2 & \text{if } n > 0 \text{ is odd} \end{cases}$$

□ For example,

- $2^4 = 2^{(4/2)^2} = (2^{4/2})^2 = (2^2)^2 = 4^2 = 16$
- $2^5 = 2^{1+(4/2)^2} = 2(2^{4/2})^2 = 2(2^2)^2 = 2(4^2) = 32$
- $2^6 = 2^{(6/2)^2} = (2^{6/2})^2 = (2^3)^2 = 8^2 = 64$
- $2^7 = 2^{1+(6/2)^2} = 2(2^{6/2})^2 = 2(2^3)^2 = 2(8^2) = 128.$

A Recursive Squaring Method

Algorithm Power(x, n):

Input: A number x and integer $n = 0$

Output: The value x^n

if $n = 0$ then

return 1

if n is odd then

$y = \text{Power}(x, (n - 1) / 2)$

return $x \cdot y \cdot y$

else

$y = \text{Power}(x, n / 2)$

return $y \cdot y$

Analyzing Recursive Squaring Method

Algorithm Power(x, n):

Input: A number x and integer $n = 0$

Output: The value x^n

if $n = 0$ then

return 1

if n is odd then

$y = \text{Power}(x, (n - 1) / 2)$

return $x \cdot y \cdot y$

else

$y = \text{Power}(x, n / 2)$

return $y \cdot y$

Each time we make a recursive call we halve the value of n ; hence, we make $\log n$ recursive calls. That is, this method runs in $\log n$ time.

It is important that we used a variable twice here rather than calling the method twice.

Binary Recursion

- Binary recursion occurs whenever there are two recursive calls for each non-base case.
- Define a recursive way to print the ticks and numbers like an English ruler:

2 inch ruler w/
major tick length 4

```

----- 0
-
--
-
---
-
--
-
----- 1
-
--
-
----
-
--
-
----- 2
  
```

(a)

```

----- 0
-
--
-
---
-
--
-
----- 1
-
--
-
----- 1
  
```

(b)

```

--- 0
-
--
-
--- 1
-
--
-
--- 2
-
--
-
--- 3
  
```

(c)

3 inch ruler w/
major tick length 3

1 inch ruler w/
major tick length 5

Recursive Method for Ruler Drawing

```
// draw one tick
public static void drawOneTick(int tickLength, int tickLabel)
{
    for (int i = 0; i < tickLength; i++)
        System.out.print("-");
    if (tickLabel >= 0)
        System.out.print(" " + tickLabel);
    System.out.print("\n");
}

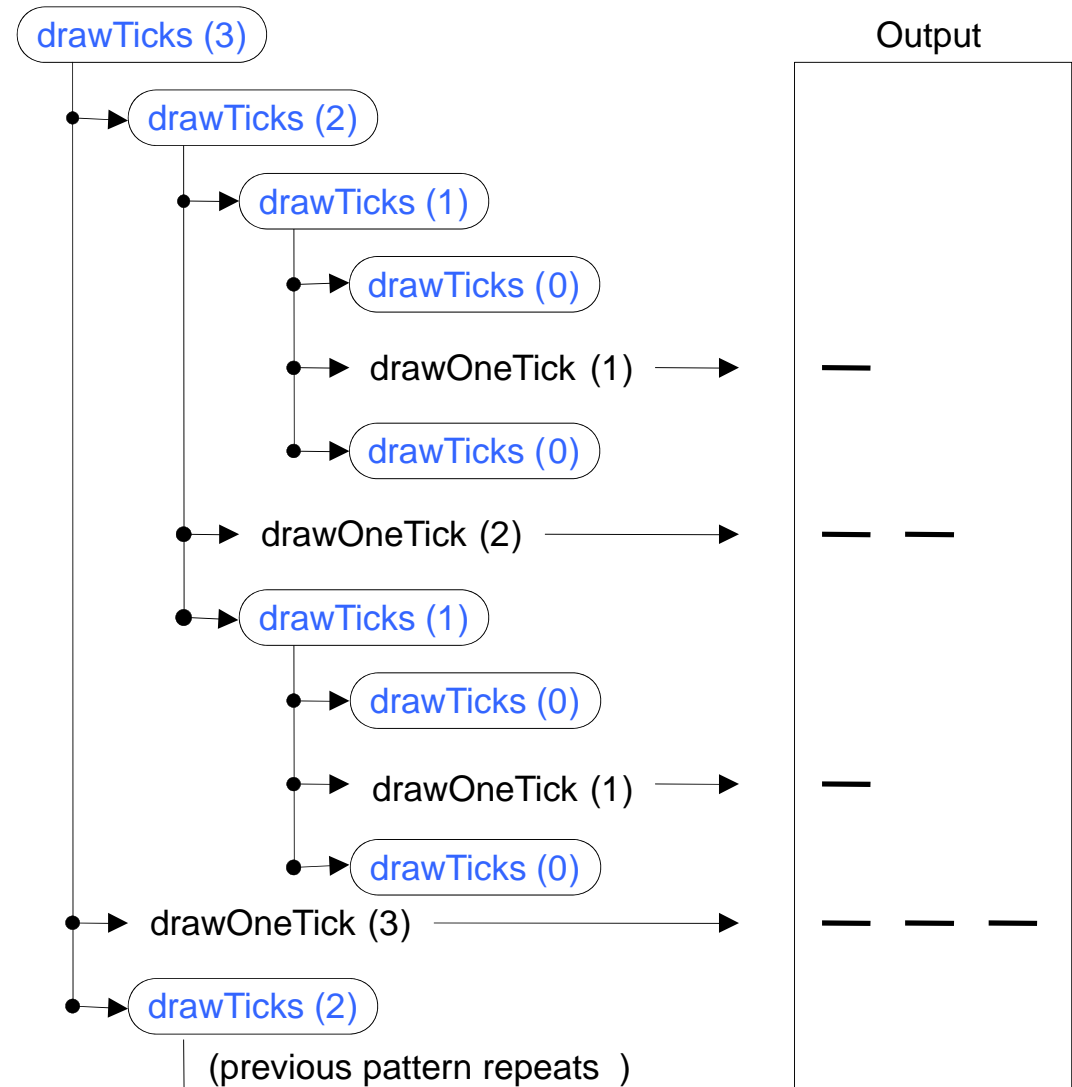
public static void drawTicks(int tickLength)
{ // draw ticks of given length
    if (tickLength > 0)
    { // stop when length drops to 0
        drawTicks(tickLength - 1); // recursively draw left ticks
        drawOneTick(tickLength); // draw center tick
        drawTicks(tickLength - 1); // recursively draw right ticks
    }
}

public static void drawRuler(int nInches, int majorLength)
{ // draw ruler
    drawOneTick(majorLength, 0); // draw tick 0 and its label
    for (int i = 1; i <= nInches; i++)
    {
        drawTicks(majorLength - 1); // draw ticks for this inch
        drawOneTick(majorLength, i); // draw tick i and its label
    }
}
```

Visualizing the DrawTicks Method

□ An interval with a central tick length $L \geq 1$ is composed of the following:

- an interval with a central tick length $L-1$,
- a single tick of length L ,
- an interval with a central tick length $L-1$.



Another Binary Recursive Method

□ Problem: add all the numbers in an integer array A :

Algorithm BinarySum(A, i, n)

Input: array A and non-negative integer indices i and n

Output: sum of n elements in A starting at index i

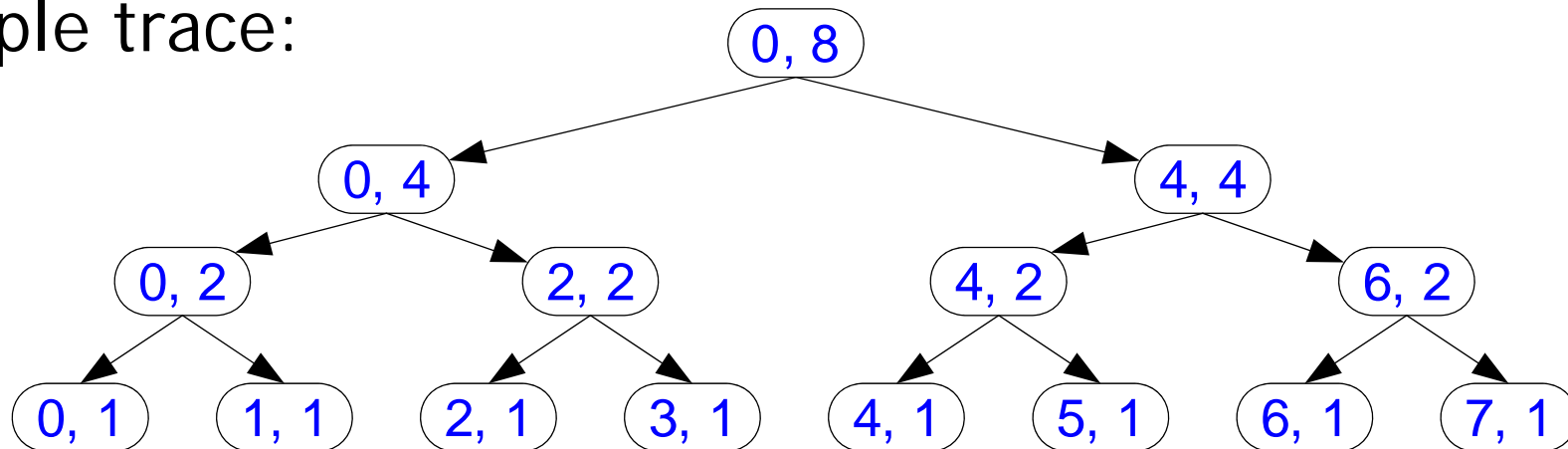
if $n = 1$ **then**

return $A[i]$

end if

return BinarySum($A, i, n/2$) + BinarySum($A, i + n/2, n/2$)

□ Example trace:



Computing Fibonacci Numbers

□ Fibonacci numbers are defined recursively:

- $F_0 = 0$
- $F_1 = 1$
- $F_i = F_{i-1} + F_{i-2}$ for $i > 1$.

□ As a recursive algorithm (first attempt):

Algorithm BinaryFib(k)

Input: non-negative integer k

Output: k -th Fibonacci number F_k

if $k \leq 1$ **then**

return k

else

return BinaryFib($k - 1$) + BinaryFib($k - 2$)

end if

Analyzing the Binary Recursion Fibonacci Algorithm

□ Let n_k denote number of recursive calls made by BinaryFib(k).

Then

- $n_0 = 1$
- $n_1 = 1$
- $n_2 = n_1 + n_0 + 1 = 1 + 1 + 1 = 3$
- $n_3 = n_2 + n_1 + 1 = 3 + 1 + 1 = 5$
- $n_4 = n_3 + n_2 + 1 = 5 + 3 + 1 = 9$
- $n_5 = n_4 + n_3 + 1 = 9 + 5 + 1 = 15$
- $n_6 = n_5 + n_4 + 1 = 15 + 9 + 1 = 25$
- $n_7 = n_6 + n_5 + 1 = 25 + 15 + 1 = 41$
- $n_8 = n_7 + n_6 + 1 = 41 + 25 + 1 = 67.$

□ Note that the value at least doubles for every other value of n_k . That is, $n_k > 2^{k/2}$. It is exponential!

A Better Fibonacci Algorithm

- Use linear recursion instead:

Algorithm LinearFib(k)

Input: non-negative integer k

Output: Pair of Fibonacci numbers (F_k, F_{k-1})

if $k \leq 1$ **then**

return $(k, 0)$

else

$(i, j) = \text{LinearFib}(k - 1)$

return $(i + j, i)$

end if

- The number of recursive calls is linear in k
 - = Runs in $O(k)$ time.

Multiple Recursion

□ Motivating example: summation puzzles

- Assign a unique digit (0, 1, ..., 9) to each letter so that
 - pot + pan = bib
 - dog + cat = pig
 - boy + girl = baby

□ A simple (but naive) solution: enumerate all possibilities (=permutations) and test each one

- Multiple recursion: makes potentially many recursive calls (not just one or two)

Algorithm for Multiple Recursion

Algorithm PuzzleSolve(k, S, U)

Input: integer k , sequence S , and set U

Output: enumeration of all k -length extensions to S
using elements in U without repetitions

for each e in U **do**

 remove e from U // e is now being used

 add e to the end of S

if $k = 1$ **then**

 Test whether S is a configuration that solves the puzzle

if S solves the puzzle **then**

return "Solution found: " S

end if

else

 PuzzleSolve($k - 1, S, U$)

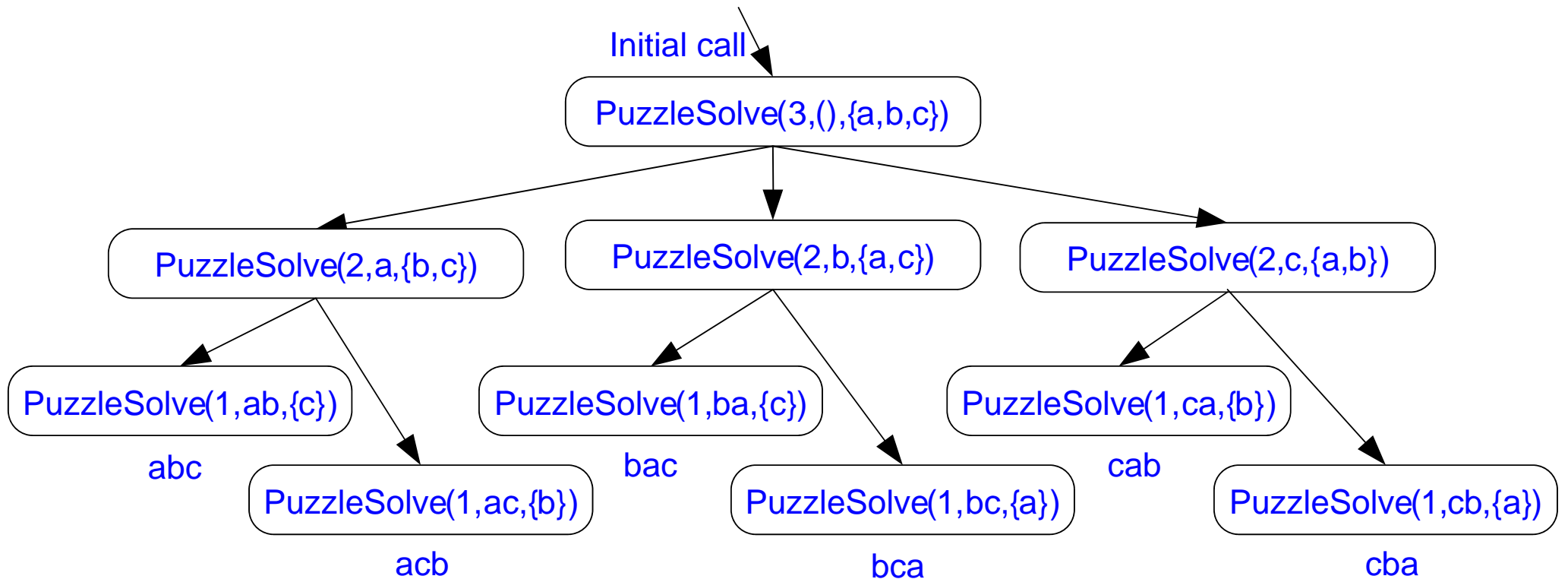
end if

 add e back to U // e is now unused

 remove e from the end of S

end for

Visualizing PuzzleSolve

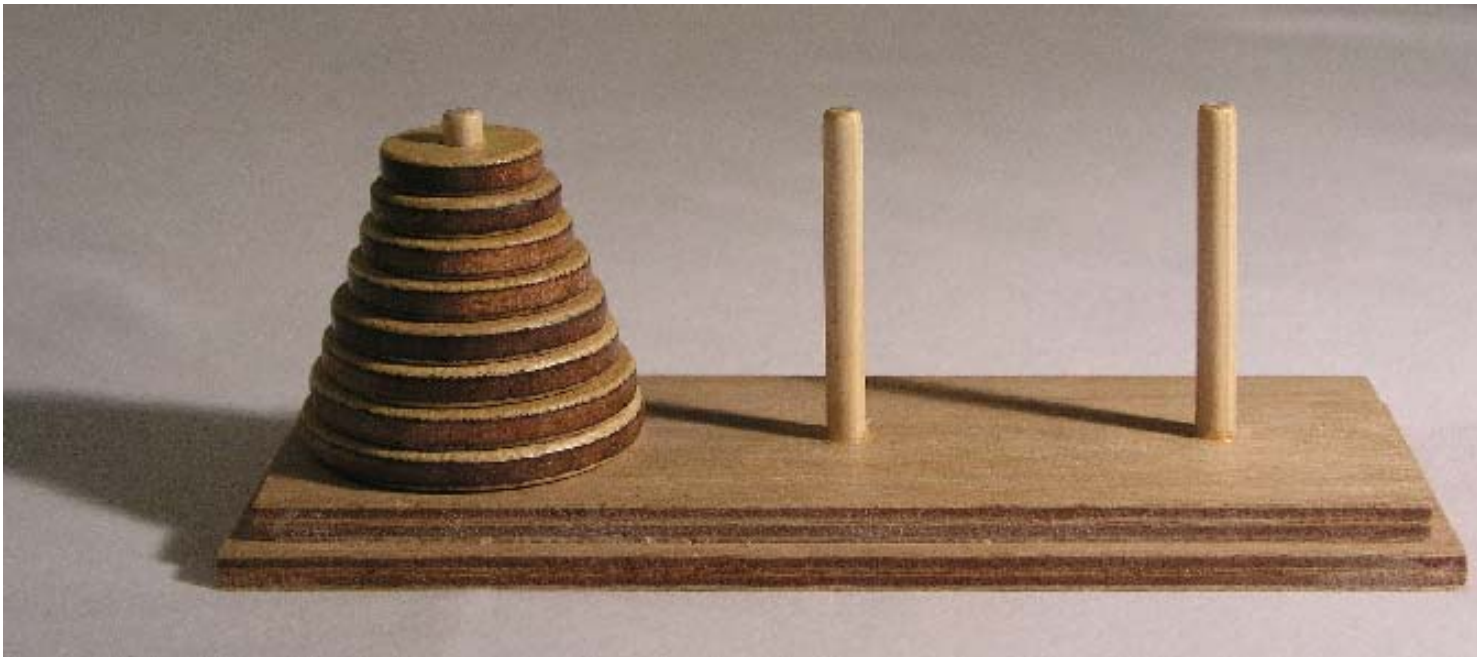


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- Can write a recursion function for addressing problems
 - Base case
 - Recursive calls
- Get to know that a program can run fast or slow depending on its algorithm

PA1

□ Implement the tower of Hanoi using recursion



Next Time

Analysis tools

HW:

- Go over the next lecture slides before the class
- Just 10 min ~ 20 min should be okay

Any questions:

- Come up with one question on what we have discussed in the class and submit at the end of the class
- 1 for typical questions or 2 for questions with thoughts or that surprised me

- Write questions at least 4 times in this semester
- You can type at KLMS